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# Internet Architectural Philosophy and the New Business Reality

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## Topics

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- ◆ architecture (as design philosophy)
- ◆ key decisions
- ◆ architecture (as reality)
- ◆ and then there is money

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## Background

- ◆ multiple unrelated efforts (early to mid 1960's)
  - packet switching theory: (Kleinrock) 1961
  - day dreaming: (Licklider's Galactic Network) 1962
  - make use of remote expensive computers: (Roberts) 1964
  - survivable infrastructure for voice and data: (Baron) 1964
- ◆ ARPANET (late 1960's)
  - Roberts ARPANET paper 1967
  - RFP for "Interface Message Processor" won by BBN 1968
  - four ARPANET hosts by 1969
  - public demo and email in 1972

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## Fundamental Goal of Internet Protocols

- ◆ multiplexed utilization of **existing** networks
  - different administrative boundaries
  - multiplexing via packets
  - networks interconnected with packet switches
    - called gateways (now called routers)
  - note: international in scope
- ◆ did not want to build a new global network
  - too expensive
  - too limiting

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# Internet Protocols Design Philosophy

- ◆ ordered set of 2nd-level goals
  - 1/ **survivability** in the face of failure
  - 2/ support **multiple types** of communications service
  - 3/ accommodate a **variety** of network types
  - 4/ permit **distributed management** of resources
  - 5/ **cost effective**
  - 6/ **low effort** to attach a host
  - 7/ **account** for use of resources
- ◆ note: no performance (QoS) or security goals
- ◆ not all goals have been met
  - management & accounting functions are limited

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# Packets!

- ◆ basic decision: use packets not circuits
  - Kleinrock's work showed packet switching to be a more efficient switching method
- ◆ packet (a.k.a. datagram)

Dest Addr	Src Addr	payload
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  - self contained
  - handled independently of preceding or following packets
  - contains destination and source **internetwork** address
  - may** contain processing hints (e.g. QoS tag)
  - no delivery guarantees**
    - net may drop, duplicate, or deliver out of order
    - reliability (where needed) is done at higher levels

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## Routing

- ◆ sub parts of the network are connected together by computers that forward packets toward destination  
these computers are called “**routers**”
- ◆ routers use destination address in packet to make forwarding decision
- ◆ routers exchange reachability information with other routers to build tables of “next hops” toward specific local networks  
exchange of reachability information done with “**routing protocol**”

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## Unreliability can be Important

- ◆ basic decision: **offer** an unreliable service
- ◆ 1st idea was to only have TCP (a reliable service)
- ◆ problems
  - not good for voice & video
    - data has to be delivered in time - retransmission for reliability causes too great a delay
  - not good for all applications
    - e.g. a debugger has to work in lossy environment
    - retransmission algorithm may vary with application
- ◆ thus: **split** IP & TCP and add UDP  
IETF just added SCTP

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## A Quote

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*“the lesson of the Internet is that **efficiency is not the primary consideration**. Ability to grow and adapt to changing requirements is the primary consideration. This makes simplicity and uniformity very precious indeed.”*

Bob Braden

## Networks as Generic

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- ◆ design requirement of working over:
  - existing networks & a wide variety of networks
- ◆ minimum set of assumptions about network
  - reasonable size packets, reasonable but not perfect delivery reliability, network-wide addressing, way to get error messages back to source, no assumption of in-order packet delivery
- ◆ “smart wires” are not much of a help
  - e.g. X.25 (reliable delivery)
  - e.g. ATM (QoS functions)
- ◆ thus it is easy to use new types of networks
  - assuming they are not too helpful (feature rich)

## *End-to-End Argument*

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- ◆ 1981 paper by Saltzer, Reed & Clark
- ◆ “smart networks” do not help
  - adding functions into network can be redundant since actual function is end-to-end
    - e.g. encryption, data reliability
  - also harder to change with new technology
    - also see Lampson *Hints for Computer System Design*
- ◆ e2e argument projected to mean
  - no per-session knowledge or state in the network
    - but some “soft-state” (auto refreshed) may be OK
  - network should be transparent to end-to-end applications

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## *Ease of Experimentation With e2e*

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- ◆ easier to experiment in an e2e environment
  - if the network is transparent then only nodes involved are the end nodes
    - note that an end node could be a 3rd party server
  - no need to get permission to experiment
- ◆ cheaper to experiment
  - can do much smaller scale experiments - down to 2 nodes than core-based services
- ◆ WWW an example of what can be done

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## Economic Driver?

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- ◆ Mark Gaynor Harvard PHD thesis
- ◆ define market uncertainty as **MU**
  - how well do you know what the customer wants
- ◆ low MU means customer wants are known
  - e.g. “voice service”
  - no opportunity to be “better” than competitor
- ◆ high MU means customer wants are not known
  - e.g. future IP-enabled voice service
  - opportunity to better match customer wants than competitor does

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## Economic Driver, contd.

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- ◆ low MU
  - commodity service
  - provide most efficient way - frequently centralized
- ◆ high MU
  - need to experiment to try to match customer want
  - note: if only one company figures it out they dominate the market
  - easier to experiment on edges
    - i.e. e2e is a innovation friendly model
  - even if its more expensive to provide service to ends

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## Smart vs. Stupid Networks

- ◆ phone network technology: self-named “Intelligent Network” (IN)
  - many network-based services
    - admission control, number translation, accounting, ...
- ◆ Isenberg’s *Rise of the Stupid Network* compared phone network’s “Intelligent Network” to Internet
  - Isenberg’s basic messages:
    - network (i.e. carrier) -based services slow to change
    - voice is not all there is
    - carrier gets in the way
    - just “deliver the bits” works

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## But!!

- ◆ a “stupid network” is a commodity service
  - the price of a commodity service is driven by the stupidest vendor
- ◆ hard to make money delivering commodity services
- ◆ new network infrastructure is very expensive
  - fiber optic cables (with installation) & hardware
- ◆ access rights can also be very expensive
  - e.g. wireless spectrum licenses
- ◆ carriers need something else to make money
  - common dream is that services or content will save the day
  - may be a false dream



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## But!! (2)

- ◆ packets w/o circuits cause problems
  - can not do guaranteed QoS
  - can not control path packets take
  - can not reserve capacity for application
  - security control harder
  - do not have logical “wire” back to source
  - management harder
  - can not see data patterns on the network
  - finding non-catastrophic failures harder
  - service provider interconnections harder
  - no clean interface
- ◆ lack of useful formal tools to describe performance

# !QoS

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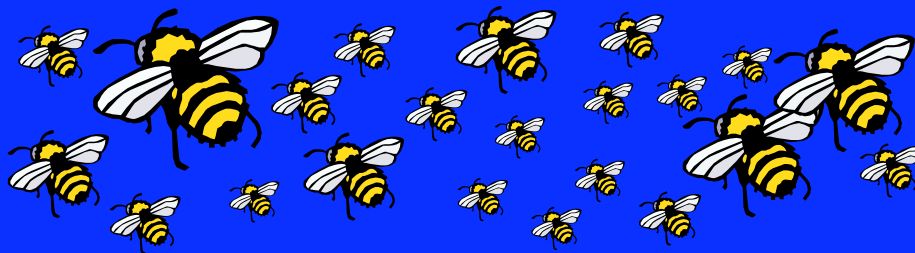
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## Conceptualization Problem

- ◆ fundamental disconnect between “Internet” and “phone” people “bell-heads vs. net-heads”
- ◆ by their definition the Internet can not work and must be fixed - they will rescue us

***“You can not build corporate network out of TCP/IP.”***

IBM circa 1992

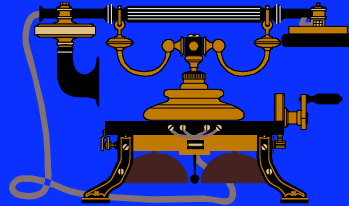


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## Traditional Phone Network

- ◆ circuits & “smart network”
- ◆ connection-oriented
- ◆ hard state in network devices
- ◆ fragile
- ◆ central resource control
- ◆ socialist? "for the good of all"
- ◆ applications in network
  - e.g., phone switch
  - end-to-end touch-tone signaling was a mistake
- ◆ predictable development path
  - extended development cycle

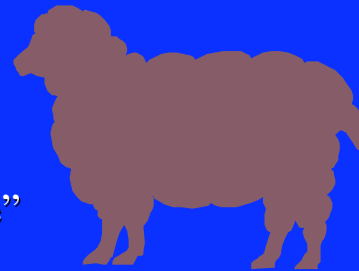


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## Internet

- ◆ packets & e2e
- ◆ soft state in network devices
- ◆ resilient
- ◆ competitive resource control
- ◆ capitalist? "individual initiative"
  - but too much selfishness hurts all
  - must play by the same rules - but no enforcement
  - the tragedy of the commons**
- ◆ applications in hosts at edges (end-to-end)
  - and in 3rd party servers anywhere on the net
- ◆ hard to predict developments
  - chaos at the rate of “Internet time”



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## More Conceptualization Problems

- ◆ service provided by 3rd parties - not only by carriers

different from phone world

- ◆ a quote from an IETF telephony mailing list

Hi Roy,

I still don't understand why it is a "users" choice where the "services" are executed - I would have thought that this would be networks choice

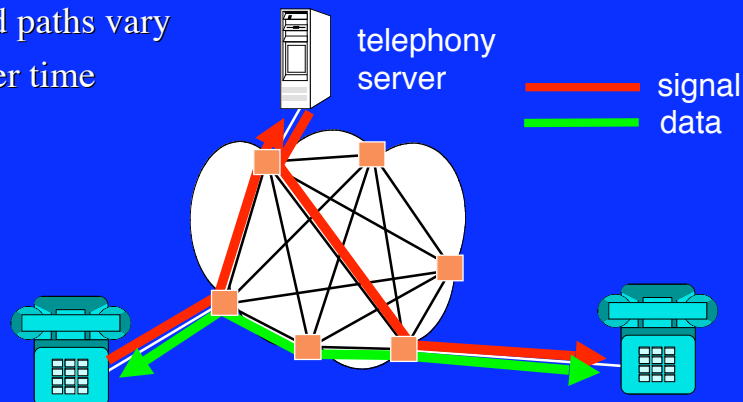
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## Disjoint Control and Data Paths

- ◆ signaling and data paths in Internet may not coincide

and paths vary  
over time

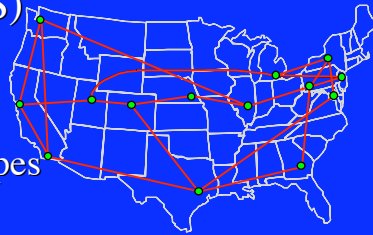


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## Circuits in the Internet

- ◆ do not seem to go away (MPLS)
- ◆ used for traffic engineering
  - city-pair pipes
  - maybe class of service city-pair pipes
- ◆ and customer connections
- ◆ finer grain (instance of application) use still pushed
- ◆ remember the fate of ATM
  - circuit - used for trunks not flows
  - QoS - ignored (ATM not end-to-end)
  - link sharing - may make sense
  - as the bearer service - did not make it



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## IP as a Common Bearer Service

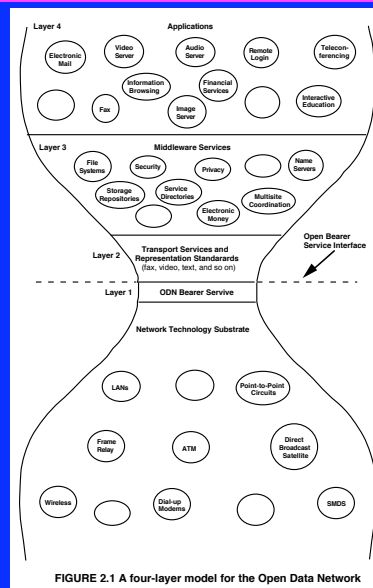


FIGURE 2.1 A four-layer model for the Open Data Network

From: Realizing the Information Future

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## Net is No Longer Transparent

- ◆ end-to-end argument says the net should be transparent
  - i.e. packet not modified in transit (other than TTL)
  - global-scope internetwork address
  - i.e., packet goes to address in destination address field
- ◆ transparency now gone in some cases
  - NATs, firewalls, proxies, content caches, TCP reshapers
  - replace addresses, intercept traffic, insert traffic
- ◆ other issues
  - wiretapping, taxation, content filtering

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## NAT/Firewall/Cache Issues

- ◆ can not trust IP address as end-to-end
  - breaks IPSec, not sure who you are talking to
- ◆ applications with addresses in data
  - have to have application-specific support (ALG) in devices
  - deploying new application requires approval of net manager
- ◆ dynamic port usage
  - ALG must snoop on application traffic
  - ALG must understand application logic
- ◆ new IETF effort to develop generic signaling
  - may help some
  - but will not make these devices transparent

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## Trust-Free Environment

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- ◆ original Internet architecture assumed a trustworthy environment
- ◆ no longer the case
  - mistrust net itself (eavesdropping, reliability etc)
  - mistrust that you are talking to the right end point
    - e.g., proxy, redirect, spoofing (MAC & IP address)
  - unsolicited correspondence (spam)
  - anonymity hard to get
  - mistrust own hardware and software
  - 3rd parties insist on being in the middle
    - filters, wiretapping, ...

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## Summary of Architectural Points

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- ◆ datagram-based network
  - not circuit switched
- ◆ network of networks
  - different parts under different management
- ◆ minimize per-session state in network
  - some auto-refreshed state is OK
- ◆ end-to-end model maximizes flexibility
  - network does not need to know what you are doing
- ◆ “smart wires” can get in the way
  - e.g., nested control loops
- ◆ reliable delivery is an option
  - not a requirement

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## Key Decisions

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- ◆ a few key decisions brought us here to the Internet of today
- ◆ but there was no way to predict where we are now
- ◆ **unplanned parenthood**

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## 10 Decisions That Made a Difference

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- ◆ support existing networks
- ◆ **datagram**-based
- ◆ creating the **router** function
- ◆ split TCP **and** IP
- ◆ DARPA fund Berkeley to add TCP/IP to **UNIX**
- ◆ CSNET and **CSNET/ARPANET** deal
- ◆ NSF **require TCP/IP** on NSFnet
- ◆ ISO **turn down** TCP/IP
- ◆ NSF Acceptable Use Policy (**AUP**)
- ◆ **minimal** regulation

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## Internet Architecture

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- ◆ #1 goal of original Internet protocols was to deal with a **network of networks**
  - not a single type of network
  - not under one management
- ◆ networks interconnected at datagram level
  - no session-aware logic at interconnections
- ◆ bi-lateral interconnection agreements
  - “customer” - buy transit service to “the Internet”
  - “peer” - cost sharing connection to a network and its customers

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## Customer Interconnection

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- ◆ one network pays another for access to “the Internet”
  - paying network can be Internet service provider (ISP) or enterprise
  - only as useful as resulting coverage
    - “Metcalfe’s Law”: value of network increases by square of the number of reachable nodes
- ◆ customer can move business to another network if they do not like the service
  - may have to renumber to preserve addressing topology

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## Peering

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- ◆ business decision
  - no current regulations
- ◆ it can be cost effective for two networks to interconnect sharing the costs of the links
  - interconnection can be at “public peering points” or using dedicated links between networks
- ◆ but only “see” other network and their customers
  - not the other network’s other peers
- ◆ must peer with all large networks to get “the Internet”
  - or be a customer to another network (or networks)

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## Multi-Homing

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- ◆ one network (ISP or enterprise) can connect to more than one other network
  - for redundancy and reliability
  - called “multi-homed”
- ◆ causes some complexity in the routing setup

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## Public Peering Points

- ◆ 3 originally designated by National Science Foundation (NSF) as part of the breakup of the NSFnet
- ◆ now many local peering points around the world but telcom costs can discourage use in some countries  
cheaper to get lines to US than within country
- ◆ level-2 interconnect  
like an local area network (e.g. an Ethernet)  
i.e. not involved in IP-level routing

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## Private Peering

- ◆ two ISPs can agree to interconnect sharing costs  
“you buy and run one line, I’ll buy and run another”  
peering list normally private
- ◆ ISPs have minimum criteria before peering will be considered  
some publish the criteria
- ◆ criteria normally include  
minimum level of interconnect traffic, traffic balance,  
backbone size, geographic scope,  
competent network operations center

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## Tier 1 ISPs

- ◆ some big ISPs are referred to as “Tier 1 ISPs”
- ◆ no real externally verifiable definition
- ◆ general concept:
  - “an ISP that gets most of not all of its connectivity from peering, not by being a transit customer”
  - i.e. a Tier 1 ISP is one that is connected to the other Tier 1 ISPs

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## Interconnection Pattern

- ◆ no explicit network hierarchy assumed
- ◆ no specific pattern to ISP interconnections
  - other than that peering tends to be between networks of the same basic size
  - but not always - can have business reasons for mismatch
- ◆ peering and transit connections can appear random
- ◆ **notes:**
  - most traffic does not flow through Tier 1 ISPs
  - many “lower-level” interconnections
  - hard (impossible) to know relative sizes of ISPs

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## Current Internet Architecture



## Money

- ◆ “but who is going to make money at that?”

John McQuillan

- ◆ how is the carrier supported?

“we do not know how to route money”

Dave Clark

- ◆ carrier wants a piece of the action

e.g., WAP, AT&T proposal

- ◆ is content king?

factoid: total US movie revenue  $\approx$  2 weeks of US phone charges

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## iMode: A Model?

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- ◆ DoCoMo's iMode service
  - more than 30 million subscribers
  - 9.6Kb data service
  - 50,000 iMode compatible sites
  - DoCoMo works with less than 10% of them
    - does billing, runs servers etc
    - rest are on their own
- ◆ key decision: open access ( NOT WAP!)
  - makes service more attractive
  - DoCoMo charges monthly fee and for data transferred

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## More on Money

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- ◆ QoS does not seem to be a useful charging base
- ◆ differentiated by application is an intelligence test
  - railroads in US used to do this (Rhode Island Line )
  - not enough will fail the test
- ◆ and then there is all that fiber
- ◆ do municipalities have a role?

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## A Bit More on Money

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- ◆ what happened to that \$ trillion anyway?
  - few infrastructures pay for themselves
  - the Internet is not an exception
- ◆ is there a difference now that the fiber is “free”?

## Last Word

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- ◆ Internet “too important to fail” (?)
  - what about ISPs (can you say “KPNQuest”?)
- ◆ will there be anyone left standing other than the telcos?
  - what can they see from their point of view?
- ◆ will you be able to say “Internet” and “business model” in the same sentence?
  - without a “no” in between

# What's Next?

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